

not display the second data unit at all and the lock E2 is not indicated to the second user until the second data unit is visible to the second user.

[0094] In FIG. 20, the second user adds a third data unit to the document and a third lock E3 around the third data unit (e.g., in the memory of the authoring application). The third data unit contains the text “Goodbye World”. The second authoring application generates (e.g., via a layer object) a working copy of the document including the third data unit and the third lock E3. The second authoring application also automatically provides (e.g., via a sync manager) the third lock E3 to the storage device. However, the second authoring application does not provide the third data unit to the storage device.

[0095] Also in FIG. 20, the first user releases the first lock E1 after the first user indicates the first user has finished editing the first data unit. For example, the first user can provide instructions to store the document when the cursor of the first user is positioned within the second data unit. In one embodiment, the second data unit may be located remote from the first data unit. In the example shown, the finalization indication provided by the first user does not indicate the first user wishes to share content changes made by the first user. Accordingly, the first authoring application does not generate (e.g., via a layer object) an upload copy containing content changes for distribution to the storage device.

[0096] In FIG. 21, the third lock E3 has been stored in the master lock table of the master copy 1250 of the document on the storage device. The first user computing device subsequently polls (e.g., via a sync manager) the storage device to check for updates. The first user computing device obtained (e.g., via a sync manager) a new download copy of the document from the storage device based on the master copy 1250. The new download copy contains the third lock E3. The first authoring application automatically instantiates (e.g., via a layer object) the third lock E3 into the working copy of the first authoring application. As noted above, the third lock E3 optionally may be displayed to the first user using the user interface of the first authoring application.

[0097] In FIG. 22, the second user releases the third lock E3 when the second user indicates the second user has finished editing the third data unit. The third lock E3 is removed from the metadata table of the second user. In the example block diagram shown in FIG. 22, the third lock E3 is removed from the metadata tables corresponding to the working copy of the second authoring application.

[0098] In the example shown in FIG. 22, the finalization indication provided by the second user indicates the second user wishes to share content changes with other users. For example, the second user may remove a cursor from the third data unit and then provide instructions to store the document. Accordingly, the second authoring application may generate an upload copy containing the content changes based on the working copy. In one embodiment, the upload copy also can be based on any content updates (e.g., in this case none) found in the download copy that had not yet been instantiated into the working copy. In the example shown, the upload copy includes the third data unit including the phrase “Goodbye World.”

[0099] In FIG. 23, the storage device obtains the third data unit from the upload copy of the second user computing device and merges the third data unit into the master copy 1250 of the storage device. The lock table associated with the master copy 1250 indicates the third lock E3 has been

released. As shown in FIG. 24, upon a successful push to the storage device, the second authoring application copies (e.g., via a layer object) the upload copy to the download copy and the base copy of the second cache 1234.

[0100] In FIG. 25, the first user computing device obtains (e.g., via a sync manager) a new download copy of the master copy 1250 of the document. The first authoring application determines (e.g., via a layer object) the new download copy has been obtained and determines the new download copy includes content changes. The authoring application automatically merges any new locks (none in this case) into the working copy of the cache 1214. The user interface of the first authoring application may indicate content updates are available for viewing.

[0101] In FIG. 26, the first user chooses to instantiate the content updates available in the cache 1214 into the working copy of the document. Accordingly, the authoring application merges the download copy stored in the cache 1214 with the working copy of the first authoring application. The third data unit is added to the working copy of the first authoring application and is displayed to the first user. The third lock E3 is released in the metadata table associated with the working copy of the first authoring application. The first user, therefore, may instruct the first authoring application to edit the third data unit.

[0102] Referring now to FIGS. 27-31, the user experience during the coauthoring session is described in detail below. FIGS. 27-31 illustrate changes to an exemplary user interface 2000 of a first authoring application as a first user and a second user collaboratively author a document 2010. The first user, Jane Doe, edits the document 2010 using the first authoring application and the second user, John Doe, edits the document 2010 using a different authoring application.

[0103] The document 2010 includes a first data unit 2012, a second data unit 2014, a third data unit 2016, and a fourth data unit 2018. The first user places a content lock 2022 around the first data unit 2012. For example, in one embodiment, the first user may move her cursor 2015 to the first data unit 2012. In another embodiment, the first user may provide instructions to perform one or more editing operations on the first data unit 2012.

[0104] In FIG. 28, the first user has added text to the first data unit 2012. The lock on the first data unit inhibits the second user from editing the first data unit 2012. In one embodiment, the first lock 2022 prevents the second user from editing the first data unit 2012. In another embodiment, the first lock 2022 warns the second user against editing the first data unit 2012, but does not bar editing by the second user. In yet another embodiment, the first lock 2022 merely indicates the first user has expressed an intention to edit the first data unit 2012.

[0105] In FIG. 29, the second user has placed a content lock 2024 about the third data unit 2016. Because the second user has chosen to not share his changes, the second lock may be visible on the user interface 2000 of the first authoring application, but any changes made to the third data unit 2016 by the second user are not visible yet. In the example shown, the second lock 2024 is indicated by the user interface 2000 by displaying a bracket indicating the scope of the second lock 2024. The second lock 2024 also can include a nametag indicating the owner of the second lock 2024. In FIG. 29, the nametag includes text identifying the second user and an icon. The icon can further identify a user by color, symbol, or other graphic.